

# Pregroups and distributional semantics

Antonin Delpeuch

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## The distributional hypothesis

*You shall know a word by the company it keeps.* (Firth)

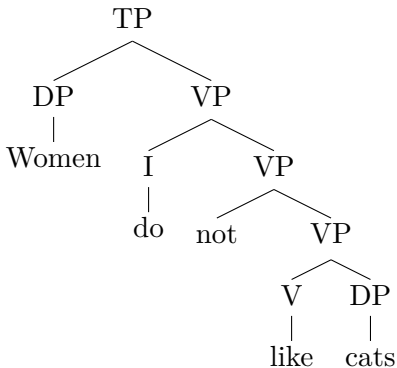
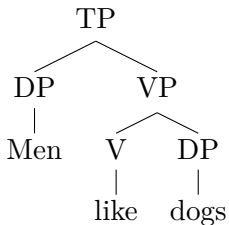
Words that occur in the same contexts tend to have similar meanings.

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- enables comparison of meanings

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In both cases, the reduction rules are formalized through a partial order.

## Definition

A pregroup is a monoid with a partial order  $\leq$  and two unary operations  $^l$  and  $^r$  such that :

- if  $a \leq b$  then  $ca \leq cb$  and  $ac \leq bc$  for all  $c$
- $x^l x \leq 1 \leq x x^l$  and  $x x^r \leq 1 \leq x^r x$

# Axioms of the Lambek calculus

## Calculus

- (a)  $z \rightarrow z$
- (b)  $(xy)z \rightarrow x(yz)$
- (c) if  $xy \rightarrow z$ , then  $x \rightarrow z/y$
- (d) if  $x \rightarrow z/y$ , then  $xy \rightarrow z$
- (e) if  $x \rightarrow y$  and  $y \rightarrow z$  then  $x \rightarrow z$

## Pregroup

reflexivity of  $\leq$

associativity of  $\cdot$

$$x = x1 \leq xyy^l = zy^l$$

$$xy = zy^l y \leq z1 = z$$

transitivity of  $\leq$

Simplified grammar of English :  $n$  for nouns,  $s$  for sentences,  $j$  for infinitive verbs.

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## Definition

A **monoidal category** is a category with a monoid law  $\otimes$  on the objects, such that for each pair of arrows ( $f : A \rightarrow B$ ,  $g : C \rightarrow D$ ) there exists  $f \otimes g : A \otimes C \rightarrow B \otimes D$  and such that

$$(f_1 \otimes g_1) \circ (f_2 \otimes g_2) = (f_1 \circ f_2) \otimes (g_1 \circ g_2)$$

In a set theoretic interpretation :  $f \otimes g$  is just the parallel evaluation of  $f$  and  $g$ .

## Definition

A **compact closed category** is a monoidal category such that for each object  $A$ , there exists :

- $A^l$  and  $A^r$
- $\eta^l : I \rightarrow A \otimes A^l$
- $\eta^r : I \rightarrow A^r \otimes A$
- $\epsilon^l : A^l \otimes A \rightarrow I$
- $\epsilon^r : A \otimes A^r \rightarrow I$

and the *yanking equalities* are satisfied.

The following equations are called *yanking equalities* :

$$(1_A \otimes \epsilon^l) \circ (\eta^l \otimes 1_A) = 1_A$$

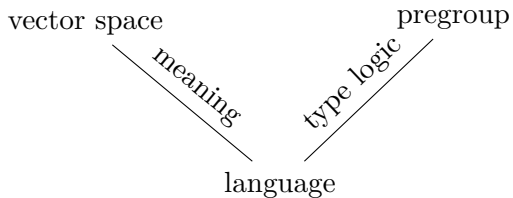
$$(\epsilon^r \otimes 1_A) \circ (1_A \otimes \eta^r) = 1_A$$

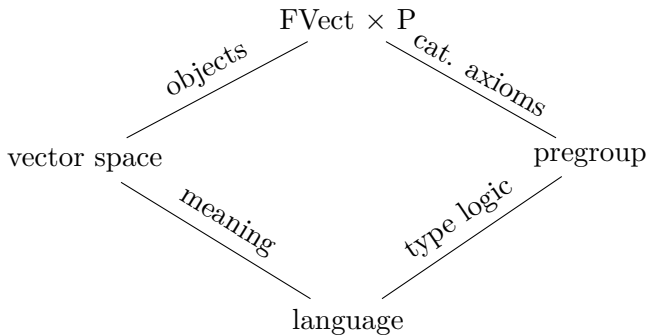
$$(\epsilon^l \otimes 1_{A^l}) \circ (1_{A^l} \otimes \eta^l) = 1_{A^l}$$

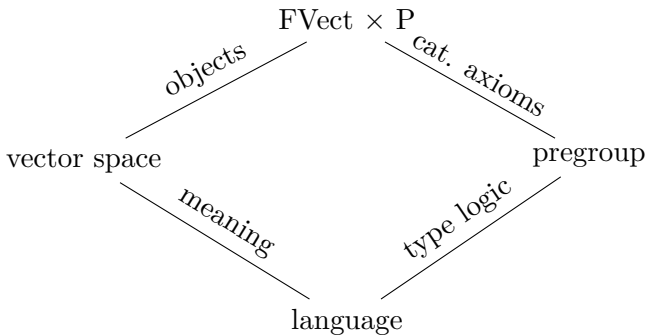
$$(1_{A^r} \otimes \epsilon^r) \circ (\eta^r \otimes 1_{A^r}) = 1_{A^r}$$

The category of vector spaces is a compact closed category :

- for the monoid law  $\otimes$  (Kronecker product) ;
- with  $A^l = A^r = A^*$  (dual) ;
- with  $\epsilon^l, \epsilon^r, \eta^l, \eta^r$  correctly defined.







$$\begin{array}{ccc}
 \text{men} & \text{like} & \text{dogs} \\
 (\overrightarrow{men} \in W, n) & (\overrightarrow{like} \in W \otimes S \otimes W, n^r sn^l) & (\overrightarrow{dogs} \in W, n)
 \end{array}$$

# Morphism associated to a reduction

If  $a \leq b$  holds, then it can be deduced from the axioms :

$$x^l x \leq 1 \leq x x^l$$

$$x x^r \leq 1 \leq x^r x$$

Similarly, in vector spaces category described before, all the morphisms can be expressed using  $\epsilon^l$ ,  $\epsilon^r$ ,  $\eta^l$  and  $\eta^r$  (and  $1_A$ , of course).

When  $a \leq b$ , we denote by  $[a \leq b]$  the morphism of vector spaces associated to this reduction.

### Meaning of a sentence

The meaning of the sentence  $(v_1 \in W_1, p_1) \otimes \dots \otimes (v_n \in W_n, p_n)$  is

$$[p_1 \cdot \dots \cdot p_n \leq s](W_1 \otimes \dots \otimes W_n)$$

.

$$S = \mathbb{R}, W = \mathbb{R}^4$$

Basis vectors for  $W$  :  $\overrightarrow{female}$ ,  $\overrightarrow{africa}$ ,  $\overrightarrow{dog}$  and  $\overrightarrow{idea}$ .

$$\begin{array}{ccc}
 \overrightarrow{women} & \overrightarrow{dogs} & \overrightarrow{like} \\
 (0.8 \quad 0.05 \quad 0.05 \quad 0.1) & (0.1 \quad 0 \quad 0.9 \quad 0) & \begin{pmatrix} 0.1 & 0.2 & 0.6 & 0.1 \\ 0.25 & 0.5 & 0 & 0.25 \\ 0.4 & 0 & 0.6 & 0 \\ 0 & 0.5 & 0 & 0.5 \end{pmatrix}
 \end{array}$$

Type reduction :  $n n^r s n^l n \leq s$ .

Morphism :  $[n n^r s n^l n \leq s] = f : W \otimes W \otimes S \otimes W \otimes W \rightarrow S$ .

$$f = \epsilon_W^r \otimes 1_S \otimes \epsilon_W^l$$

Meaning :

$$\begin{aligned} & (\epsilon_W^r \otimes 1_S \otimes \epsilon_W^l)(\overrightarrow{womeñ} \otimes \overrightarrow{like} \otimes \overrightarrow{dogs}) \\ &= (\epsilon_W^r \otimes 1_S \otimes \epsilon_W^l)\left(\sum_{i,j} c_{ij} \overrightarrow{womeñ} \otimes \overrightarrow{u}_i \otimes \overrightarrow{v}_j \otimes \overrightarrow{dogs}\right) \\ &= \sum_{i,j} c_{ij} \langle \overrightarrow{womeñ} | \overrightarrow{u}_i \rangle \langle \overrightarrow{u}_j | \overrightarrow{dogs} \rangle \\ &= 0.47025 \end{aligned}$$



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